

## BACKGROUND & MOTIVATION

Measurements of packet delay, packet loss, and throughput are widely used for broad assessment of internet performance characteristics. A study in 2000-01 by Zhang et al. [ZPDS01] evaluated the *constancy* of internet performance (round-trip delay, packet loss, and throughput) along three dimensions:

- **Mathematical constancy:** the measurements can be described by a time-invariant model.
- **Operational constancy:** performance characteristics are stable and conducive to good application performance.
- **Predictive constancy:** performance characteristics can be predicted with good accuracy.

Internet topology has undergone significant evolution since the ZPDS01 study:

- Logical connectivity is much less hierarchical
- Recent transport protocols are more latency-sensitive, e.g. BBR congestion control.
- The web has transitioned from largely static pages to supporting interactive applications
- Real-time applications are extremely latency-sensitive, e.g., games, videoconferencing.

In this study we seek to reproduce the original study, focusing first on round-trip delays, since latency has an especially significant impact on the performance of internet applications.

The cumulative distribution function (CDF) plots below visualize the cumulative likelihood of occurrence for periods of constancy in the connection between any pairs of anchors.

- About 10% of anchor-to-anchor pairs exhibit constancy on timescales of 24 hours or longer
- Roughly 50% of constancy periods last around 2-3 hours
- On closer look, the duration of constancy at several timestamps is highly variable among anchor-to-anchor measurements

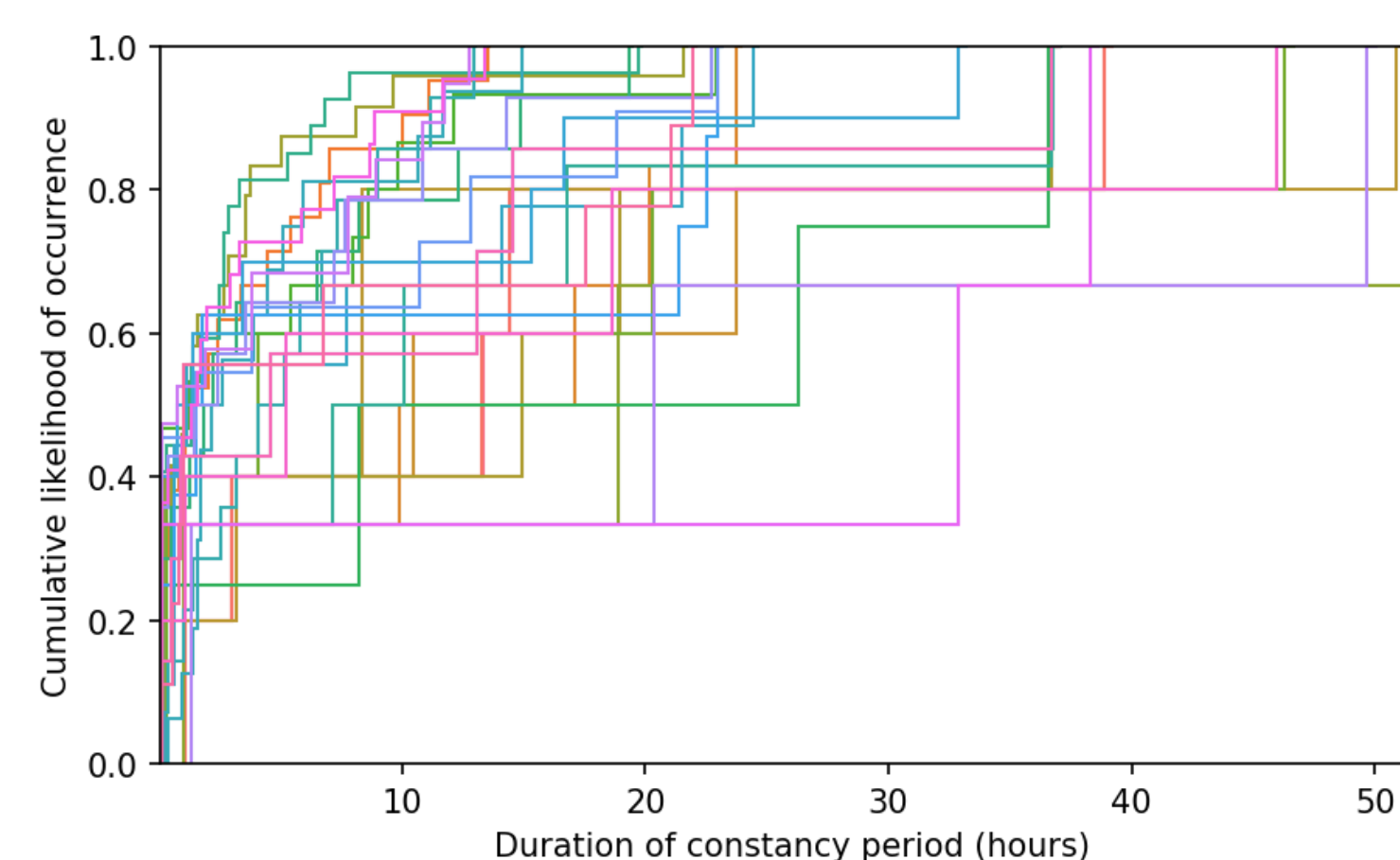


Figure 4. CDF plot of the constancy durations of each anchor-to-anchor measurement

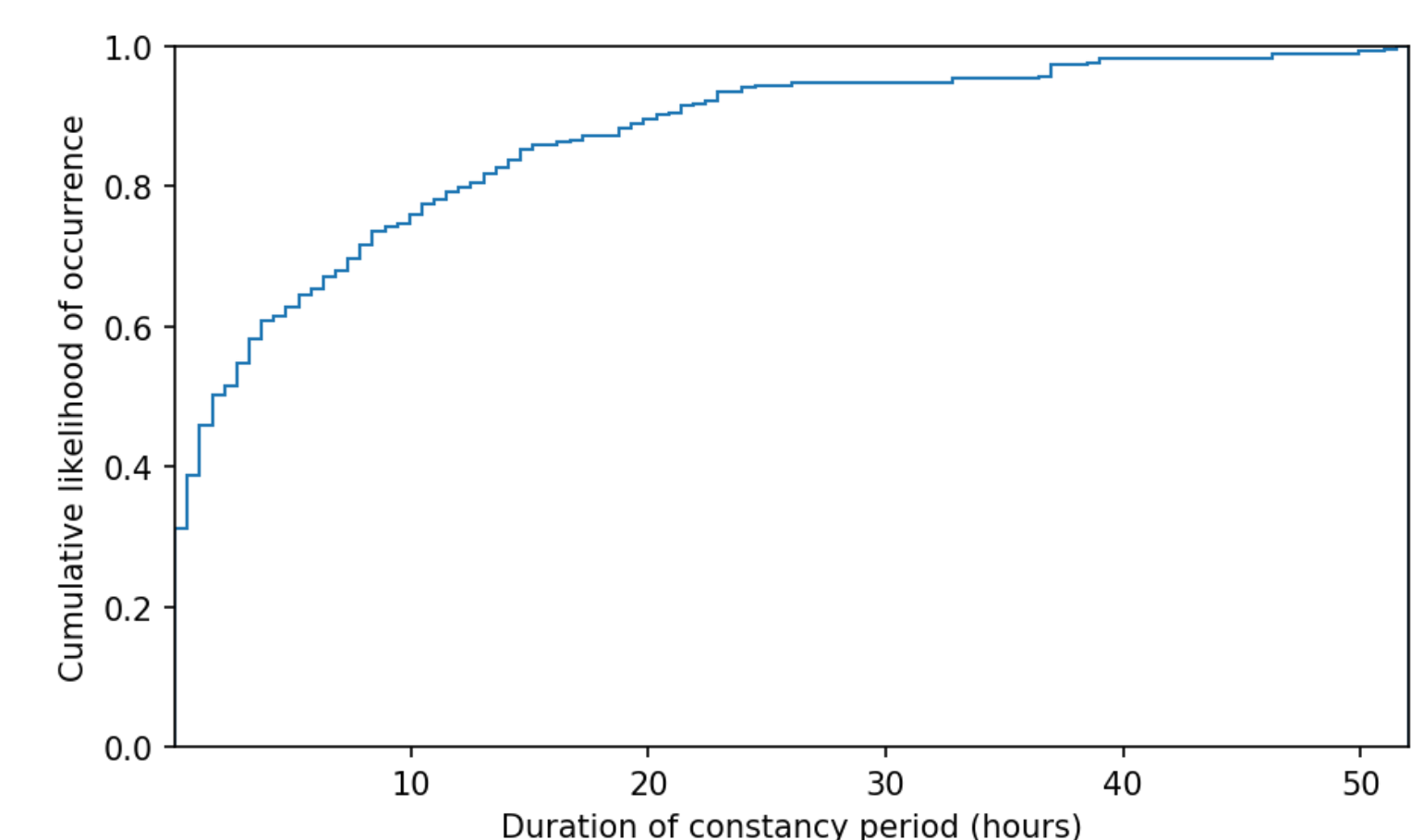


Figure 5. Aggregated CDF plot for all anchors

## DATA

We use anchor mesh data generated by the RIPE Atlas platform (<https://atlas.ripe.net>). We initially focus on RIPE Atlas anchor data because of (likely) better measurement quality.



Figure 1. Active anchors in the RIPE-Atlas network

What we collected:

- Ping measurement round-trip times (RTTs) of probe packets sent between RIPE Atlas anchors. Probes are typically sent every 4 minutes.
- First-hop latency measurements to the first router along the path from an anchor. We analyzed these to select anchors with low-variability local networks in the same spirit as Sommers et al., 2017.

We selected one anchor from each continent. The map above shows the network of active anchors.

- Almaty, Kazakhstan (AS21299)
- Palo Alto, United States (AS1280)
- Frankfurt, Germany (AS33808)
- Belo Horizonte, Brazil (AS10417)
- Johannesburg, South Africa (AS10474)
- Melbourne, Australia (AS3879)

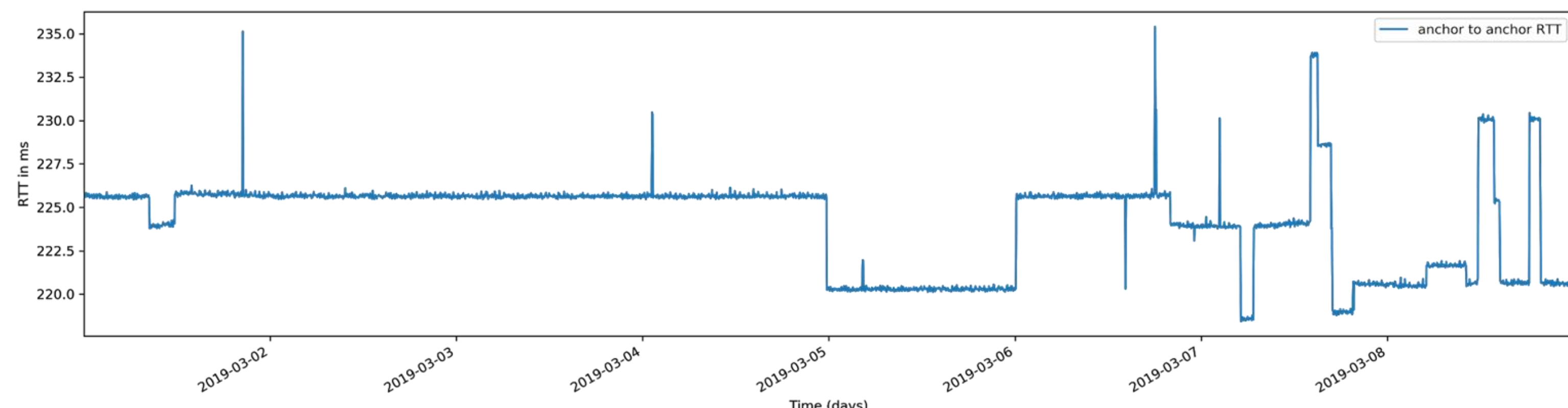


Figure 2. Ping measurement results between anchors in Belo Horizonte and Frankfurt collected between 2019/03/01 and 2019/03/09

## MATHEMATICAL CONSTANCY ANALYSIS

- We developed implementations (in Python) of the two changepoint detection algorithms (CP\_bootstrap and CP\_rank). We also used the changepoint package in R as a point of comparison (cp.var with PELT in particular).
- Using the changepoints identified, we determine periods of mathematical constancy in the data collected.

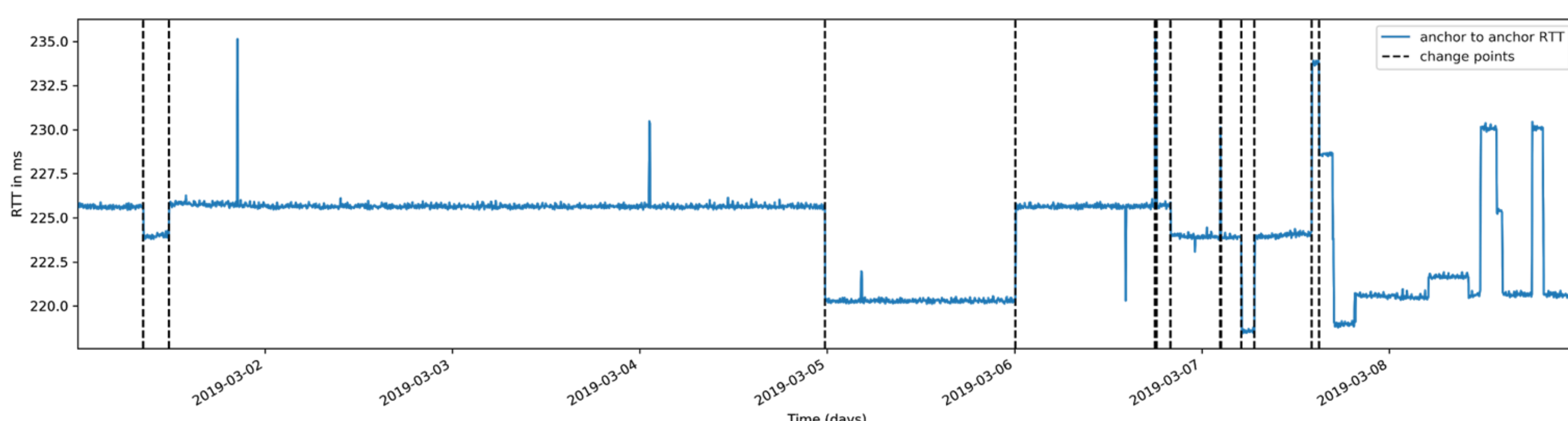


Figure 3. The vertical lines are detected changepoints. Each period between the changepoints can be described with a time-invariant model - they're mathematically "constant"

## OPERATIONAL CONSTANCY ANALYSIS

We perform a longitudinal analysis, selecting a period of two weeks every four months in the data. For the BH-Frankfurt anchor pair, the plot on the right indicates that RTTs between these two anchors generally remain within 200 - 300ms, suggesting that there is not a notable change in performance on this network. Therefore, this path is quite operationally constant.

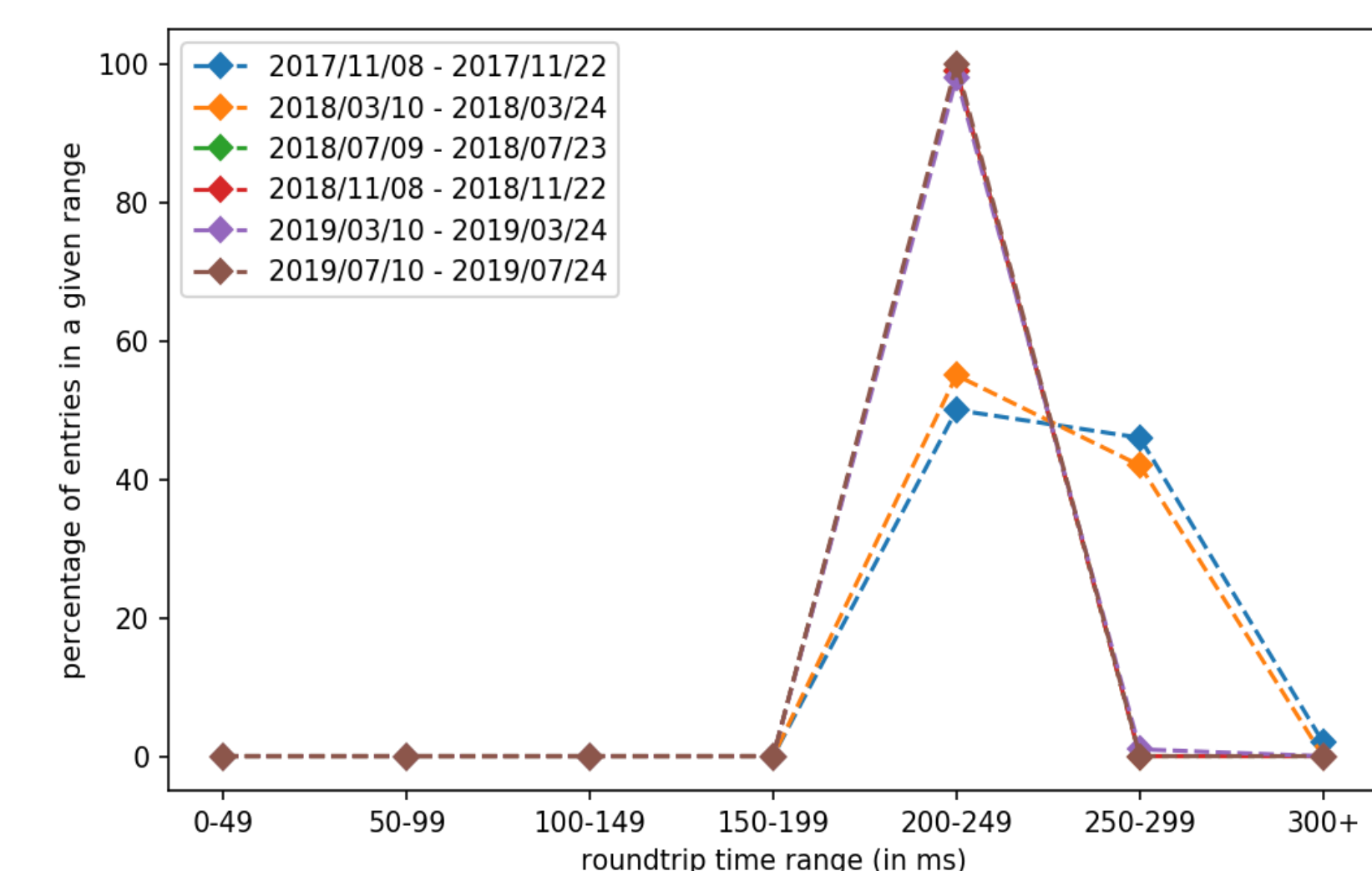


Figure 6. Longitudinal analysis of operational constancy for Belo Horizonte-Frankfurt

## RESULTS

Overall we observe:

- A wide range of CFR durations across different paths
- Unlike ZPDS01, delays along many paths may be considered operationally constant.
- No clear difference between IPv4 and IPv6 (weak evidence for longer CFRs on IPv6).
- CP\_bootstrap method is suggestive of longer CFRs, cp.var suggestive of shorter CFRs.
- Changepoint detection method matters (a lot!)

## NEXT STEPS

- Investigation of better changepoint detection algorithm(s).
- Expanding scope of data analysis.
- Additional analyses within categories of anchors, e.g., within the same continent, within tier-1 providers.
- Consideration of longer time duration.
- Longitudinal analysis of constancy (to the extent possible with available data).
- Expand analyses to consider Atlas probes (v3 and v4) more broadly.
- Analysis of "lossy" vs. "non-lossy" paths as with ZPDS01.
- Examination of constancy of packet loss and throughput.

## REFERENCES

- Zhang Y., Duffield N., Paxson V., and Shenker S. On the constancy of Internet Path Properties. Proceedings of IMW, 2001.
- Sommers J., Durairajan R., and Barford P. Automatic metadata generation for active measurement. Proceedings of IMC, 2017.
- Thanks to the RIPE Atlas team for making measurement data available to the research community.